

# AM33: The FleetTwo Dock

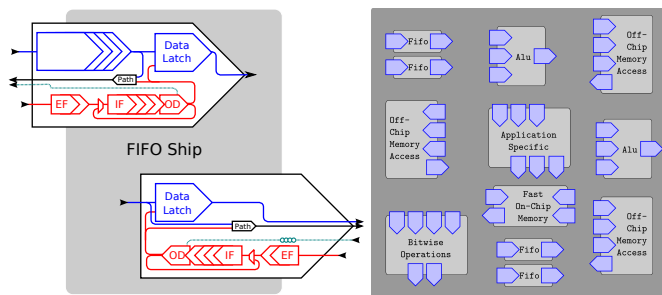
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## Abstract

### Changes:

- 16-Jun **When a torpedo strikes, ILC is set to 1**  
**Only move can be torpedoed (removed I-bit from set/shift)**
- 11-Jun Changed all uses of "Payload" to "Immediate" (not in red)  
Reworked encoding of `set` instruction
- 06-Jun Factored in Russell Kao's comments (thanks!)  
Added mechanism for setting C-flag from fabric even on outboxes
- 05-Jun Made OLC test a predicate-controlled condition  
Rewrote "on deck" section  
Added "unset" value for ILC
- 30-Apr Changed DP to `DataPredecessor` for clarity  
added comment about address-to-path ship  
changed DST field of `set` instruction from 2 bits to 3  
changed the order of instructions in the encoding map
- 23-Apr added epilogue fifo to diagrams  
indicated that a token sent to the instruction port is treated as a torpedo
- 18-Apr replaced `setInner`, `setOuter`, `setFlags` with unified `set` instruction  
replaced `literal` with `shift` instruction
- 17-Apr Made all instructions except `setOuter` depend on `OLC > 0`  
Removed ability to manually set the C flag  
Expanded predicate field to three bits  
New literals scheme (via shifting)  
Instruction encoding changes made at Ivan's request (for layout purposes)  
Added summary of instruction encodings on last page



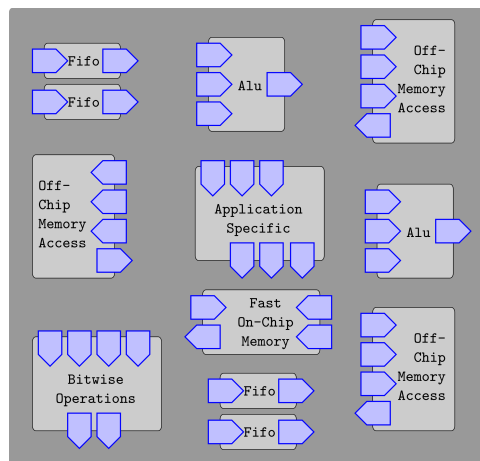
## 1 Overview of Fleet

A Fleet processor is organized around a *switch fabric*, which is a packet-switched network with reliable in-order delivery. The switch fabric is used to carry data between different functional units, called *ships*. Each ship is connected to the switch fabric by one or more programmable elements known as *docks*.

A *path* specifies a route through the switch fabric from a particular *source* to a particular *destination*. The combination of a path and a single word to be delivered is called a *packet*. The switch fabric carries packets from their sources to their destinations. Each dock has two destinations: one for *instructions* and one for *data*. A Fleet is programmed by depositing instruction packets into the switch fabric with paths that will lead them to instruction destinations of the docks at which they are to execute.

When a packet arrives at the instruction destination of a dock, it is enqueued for execution. Before the instruction executes, it may cause the dock to wait for a packet to arrive at the dock's data destination or for a value to be presented by the ship. When an instruction executes it may consume this data and may present a data value to the ship or transmit a packet.

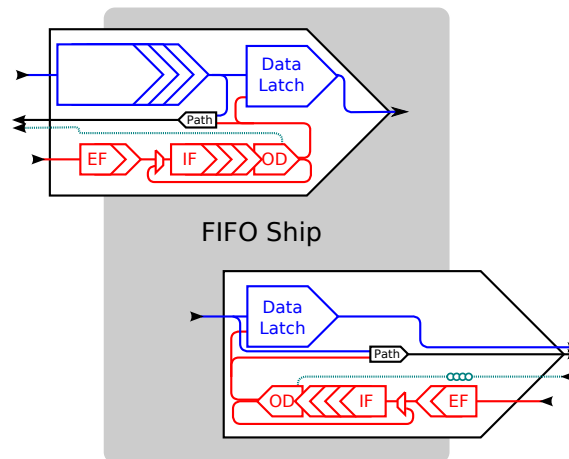
When an instruction sends a packet into the switch fabric, it may specify that the payload of the packet is irrelevant. Such packets are known as *tokens*, and consume less energy than data packets.



Overview of a Fleet processor; dark gray shading represents the switch fabric, ships are shown in light gray, and docks are shown in blue.

## 2 The FleetTwo Dock

The diagram below represents a conceptual view of the interface between ships and the switch fabric; actual implementation circuitry may differ.



An “input” dock and “output” dock connected to a ship. Solid blue lines carry either tokens or data words, red lines carry either instructions or torpedoes, and dashed lines carry only tokens.

Each dock consists of a *data latch*, which is as wide as a single machine word and a *pump*, which is a circular fifo of instruction-width latches. The values in the pump control the data latch. The dock also includes a *path latch*, which stores the path along which outgoing packets will be sent.

Note that the pump in each dock has a destination of its own; this is the *instruction destination* mentioned in the previous section.

From any source to any dock’s data destination there are two distinct paths which differ by a single bit. This bit is known as the “signal” bit, and the routing of a packet is not affected by it; the signal bit is used to pass control values between docks. Note that paths terminating at an *instruction destination* need not have a signal bit.

### 3 Instructions

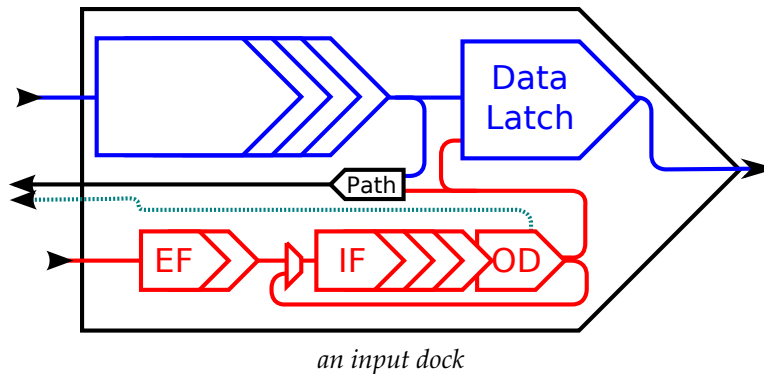
In order to cause an instruction to execute, the programmer must first arrange for that instruction word to arrive in the data latch of some output dock. For example, this might be the “data read” output dock of the memory access ship or the output of a fifo ship. Once an instruction has arrived at this output dock, it is *dispatched* by sending it to the *instruction port* of the dock at which it is to execute.

Each instruction is 26 bits long, which makes it possible for an instruction and an 11-bit path to fit in a single word of memory. This path is the path from the *dispatching* dock to the *executing* dock.



#### 3.1 Life Cycle of an Instruction

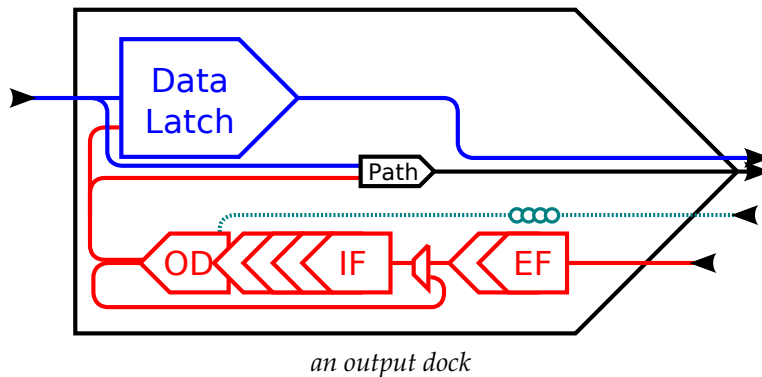
The diagram below shows an input dock for purposes of illustration:



Note the mux on the path between EF (epilogue fifo) and IF (instruction fifo); this is known as “the hatch”. The hatch has two states: sealed and unsealed. When the machine powers up, the hatch is unsealed; it is sealed by the tail instruction and unsealed whenever the outer loop counter is set to zero (for any reason<sup>1</sup>).

When an instruction arrives at the epilogue fifo (EF), it waits there until the hatch is in the unsealed state; the instruction then enters the instruction fifo. When an instruction emerges from the instruction fifo, it arrives at the “on deck” (OD) stage, where it may execute.

<sup>1</sup>this includes 0LC being decremented to zero, a set instruction, or the occurrence of a torpedo



### 3.1.1 Torpedoes

A token sent to an instruction destination is called a *torpedo*. When a torpedo arrives at the tail of EF, it is deposited in a waiting area (not shown) rather than being enqueued into EF.

There is a latch (not shown) called the *torpedo acknowledgment path latch* (TAPL) which stores a path. When a torpedo is consumed (see section "On Deck"), a token is sent along the path held in this latch.

## 3.2 Format of an Instruction

All instruction words have the following format:



- The I bit stands for *Interruptible*, and indicates if an instruction is vulnerable to torpedoes. This bit only appears in move instructions.
- The OS ("One Shot") bit indicates whether or not this instruction can pass through the pump more than once. If set to 1, then the instruction is a "one-shot" instruction, and does not pass through the instruction fifo more than once.
- The P bits are a *predicate*; this holds a code which indicates if the instruction should be executed or ignored depending on the state of flags in the dock.

### 3.3 Loop Counters

A programmer can perform two types of loops: *inner* loops of only one instruction and *outer* loops of multiple instructions. Inner loops may be nested within an outer loop, but no other nesting of loops is allowed.

The dock has two loop counters, one for each kind of loop:

- OLC is the Outer Loop Counter
- ILC is the Inner Loop Counter

The OLC applies to all instructions and can hold integers 0 . . MAX\_OLC.

The ILC applies only to move instructions and can hold integers 0 . . MAX\_ILC as well as a special value:  $\infty$ . When ILC=0 the next move instruction executes zero times (ie is ignored). When ILC= $\infty$  the next move instruction executes until interrupted by a torpedo. After every move instruction the ILC is reset to 1 (note that it is reset to 1, *not to 0*).

### 3.4 Flags and Predication

The pump has three flags: A, B, and C.

- The A and B flags are general-purpose flags which may be set and cleared by the programmer.
- The C flag is known as the *control* flag, and may be set by the move instruction based on information from the ship or from an inbound packet. See the move instruction for further details.

The P field specifies a three-bit *predicate*. The predicate determines which conditions must be true in order for the instruction to execute; if it is not executed, it is simply *ignored*. The table below shows what conditions must be true in order for an instruction to execute:

Code	Execute if
000:	OLC $\neq$ 0 and A=0
001:	OLC $\neq$ 0 and A=1
010:	OLC $\neq$ 0 and B=0
011:	OLC $\neq$ 0 and B=1
100:	OLC $\neq$ 0 and C=0
101:	OLC $\neq$ 0 and C=1
110:	OLC $\neq$ 0
111:	always

### 3.5 On Deck

When an instruction arrives on deck, two concurrent processes are started. No subsequent instruction may come on deck until both processes have completed:

1. Requeueing:
  - If the outer loop counter is zero ( $OLC=0$ ) or the instruction on deck is a one-shot instruction ( $OS=1$ ), do nothing.
  - *Otherwise* wait for the hatch to be sealed and enqueue a copy of the instruction currently on deck.
2. Execution:
  - If the instruction's predicate condition is not met (see section on predicates), do nothing.
  - *Otherwise* if the instruction is interruptible ( $I=0$ ) and a torpedo is present in the waiting area: consume the torpedo, set the outer loop counter to zero ( $OLC=0$ ), **set the inner loop counter to one ( $ILC=1$ )**, unseal the hatch, and transmit a token along in the *torpedo acknowledgment path latch* (TAPL).
  - *Otherwise* if  $ILC \neq 0$  or the instruction is *not* a move: execute the instruction.



## 4.2 set

The set command is used to set or decrement the inner loop counter, outer loop counter, and data latch.

		26	25	24	23	22	21	20	Dest												
	OS			P		1	0														
Immediate → OLC		19	18	17	16	15	14	13	12							6	1				
		10000				100										Immediate					
Data Latch → OLC		19	18	17	16	15	14	13	12												
		10000				010															
OLC-1 → OLC		19	18	17	16	15	14	13	12												
		10000				001															
Immediate → ILC		19	18	17	16	15	14	13	12							7	6	1			
		01000				100										0	Immediate				
$\infty$ → ILC		19	18	17	16	15	14	13	12							7					
		01000				100										1					
Data Latch → ILC		19	18	17	16	15	14	13	12												
		01000				010															
0-Extended Immediate → Data Latch		19	18	17	16	15	14	13											1		
		00100				0										Immediate					
1-Extended Immediate → Data Latch		19	18	17	16	15	14	13											1		
		00100				1										Immediate					
Update Flags		19	18	17	16	15	12			7			6	1							
		00010							nextA			nextB									
Immediate → TAPL		19	18	17	16	15	13										1				
		00001							Immediate												

The FleetTwo implementation is likely to have an unarchitected “literal latch” at the on deck (OD) stage, which is loaded with the possibly-extended literal *at the time that the set instruction comes on deck*. This latch is then copied into the data latch when a set Data Latch instruction executes.

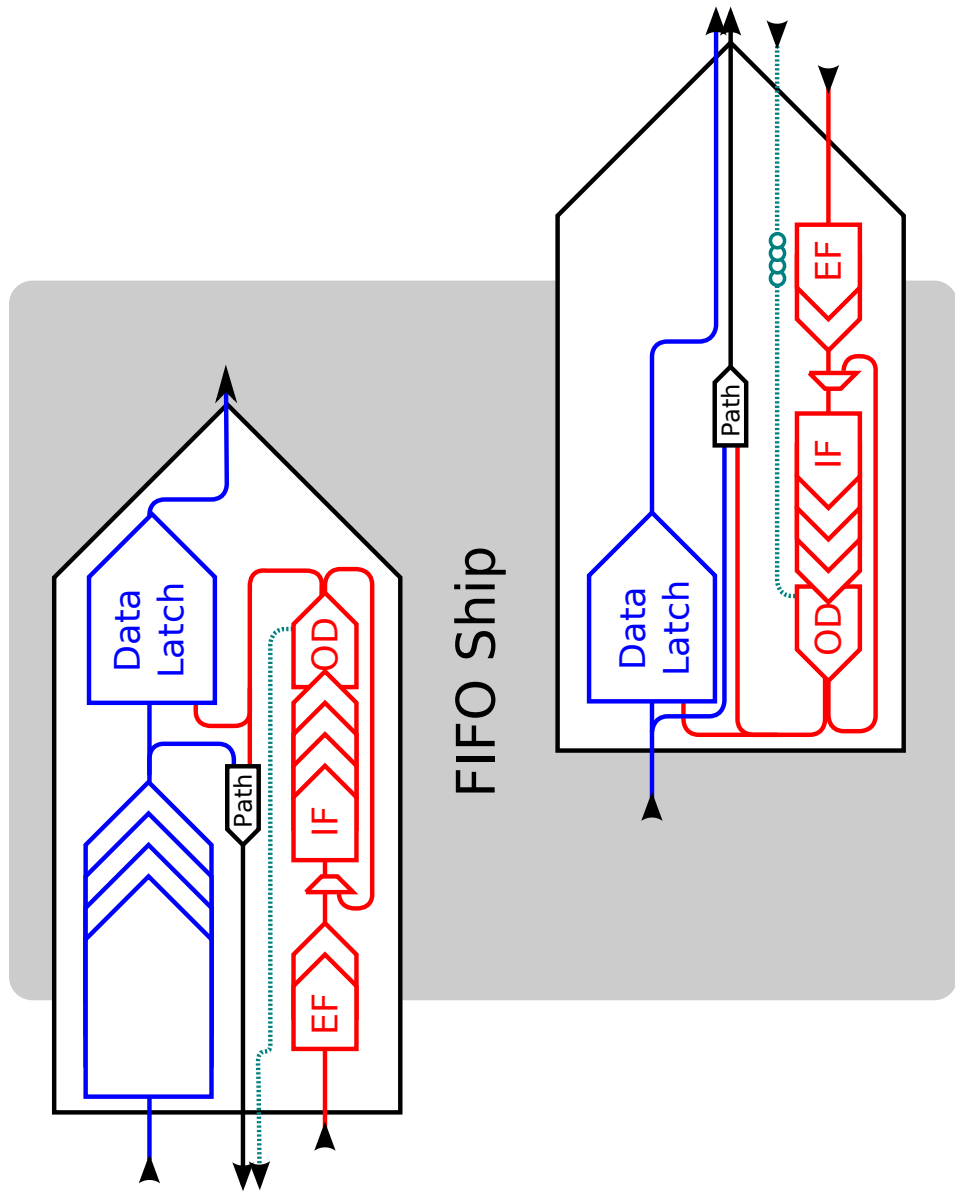
Each of the nextA and nextB fields has the following structure, and indicates which old flag values should be logically ORed together to produce the new flag value:

6	5	4	3	2	1
A	$\bar{A}$	B	$\bar{B}$	C	$\bar{C}$

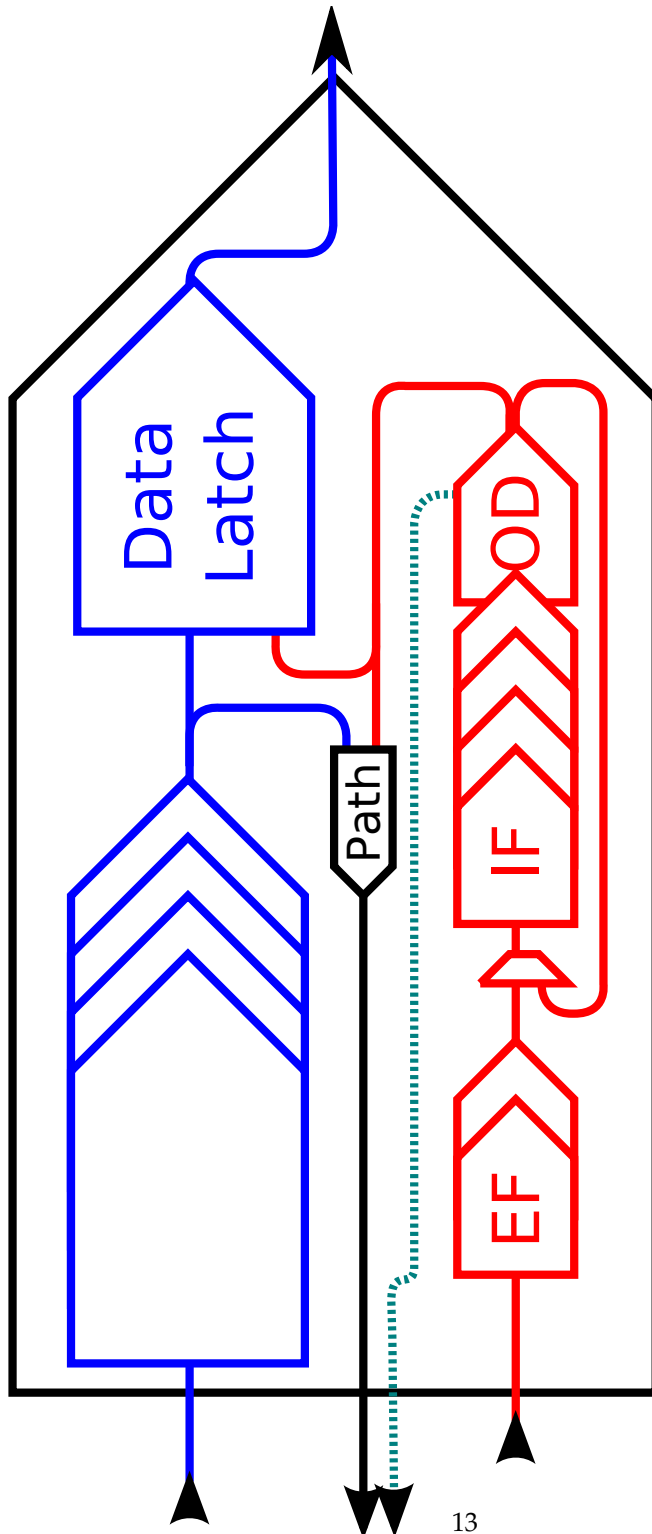
Each bit corresponds to one possible input; all inputs whose bits are set are ORed together, and the resulting value is assigned to the flag. Note that if none of the bits are set, the value assigned is zero. Note also that it is possible to produce a 1 by ORing any flag with its complement, and that set Flags can be used to create a nop (no-op) by setting each flag to itself.







Input Dock



## Output Dock

